# Guión de la exposición





- Motivación
- Índices invertidos
- Wavelet trees
- Wavelet trees sobre códigos densos WTDC
- Arrays de sufijos
- Arrays de sufijos comprimidos
- Otros índices y trabajo futuro

• ETDC + [compresor | self-index]



# Codificaciones orientadas a palabras Otros Usos:: DCC'08



### Fariña, A; Navarro, G. y Parama, J. *Word-based Statistical Compressors as Natural Language Compression Boosters*. Data compression conference. Snowbird, UT. 2008.

Introduction Related work Text preprocessing General purpose compression Boosting compression Boosting indexing Experimental results Conclusions

**Outline** 

### Introduction



- We show that most of the state-of-the-art compressors (bzip2, those from the Ziv-Lempel family and the predictive PPM-based ones) **improve** their performance if:
- They compress not the original text, but its compressed representation obtained by a word-based byte-oriented statistical compressor.

### Example:

- 1. Using End-Tagged Dense Code (ETDC) as a preprocessing step,
- 2. and then applying others (PPM,...)



Better performance (Compression ratio, compression speed and decompression speed than



### Introduction





- It also improves text indexing.
  - Text compression has been recently integrated with text indexing.
  - Self-Indexes: It is possible to construct an index which takes space proportional to the compressed text, replaces it, and permits fast indexed searching on it.
  - Examples:
    - Succinct Suffix Array (SSA) and
    - Alphabet-Friendly FM-index (AF-FMindex)

### Introduction







A self-index on the preprocessed text is **smaller** and **faster** for searching than if applied directly on T

### Outline





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### Related work





- There exist several works based on performing some text preprocessing before applying generalpurpose compressors.
  - *Mppm* from *Adiego and de la Fuente* 
    - 1<sup>st</sup> Substitutes each original word with a 2-byte id.
    - <sup>2nd</sup> Applies PPM.
  - Word replacing transformation (Skibiński, et al., 2005)
    replace original words by codewords, which index a static dictionary (in addition to other transformations) + ppm.







- There are also some works based on building a self-index over compressed text.
- WFM-index (Ferragina, 2006) builds a FM-index onto a text compressed with Tagged Huffman.
- A simple <u>alphabet-independent</u> FM-Index (Grabowski, et al., 2006) first applies a Huffman-compression and then a Burrows-Wheeler transform over it. The resulting structure can be regarded as an FM-index built over a binary sequence.

# Outline





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  - Text preprocessing
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# Related work Text preprocessing







## Semistatic compression





### Statistical semistatic compression

- Association between source symbol  $\leftarrow \rightarrow$  codeword does not change across the text.
- Direct search is possible.
- ETDC, TH son posibles

### Tagged Huffman:

•Worse compression ratio (around 35%)

•<u>Suffix-free</u>!!!



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### General purpose compression





- As a PPM compressor we chose ppmdi.
  - Uses a k-order modeler and a arithmetic encoder.
- Bzip2
  - Combines BWT, move-to-front, RLE, Huffman.
- Ziv-Lempel
  - Gzip

### Outline





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- The byte values obtained by compressing a text T with a word-based <u>byte-oriented</u> compressor shows that their frequencies are far from uniform.
  - The output of a word-based **arithmetic** <u>bit-oriented</u> compressor displays a rather homogeneous distribution.





- This idea led us to consider that the compressed file ETDC(T) (or TH(T)) was still compressible with a charbased bit-oriented compressor.
- However, this could not be a zero-order compressor, because the zero-order entropy  $(H_0)$  of ETDC(T) is too high (around 7 bpc).
- Instead, a deeper study of *k*-order entropy ( $H_k$ ) of both *T* and *ETDC(T)* exposed some interesting properties of *ETDC*.
  - A k-order modeler gathers statistics of each symbol c<sub>i</sub> by looking at the k symbols that precede c<sub>i</sub>



#### Text approx.50 Mbytes

| Plain Text |                |           |              |       |                | Text compressed with ETDC |   |                |            |    |                |      |         |
|------------|----------------|-----------|--------------|-------|----------------|---------------------------|---|----------------|------------|----|----------------|------|---------|
| k          | H <sub>k</sub> | contexts  | k            |       | H <sub>k</sub> | contexts                  | k | H <sub>k</sub> | contexts   | к  | H <sub>k</sub> | c    | ontexts |
| 0          | 4.888          | 1         | 8            | 0.972 |                | 6,345,025                 | 0 | 7.137          | 1          | 8  | 0.132          | 12,5 | 531,512 |
| 1          | 3.591          | 96        | 9            | 0.837 |                | 9,312,075                 | 1 | 6.190          | 256        | 9  | 0.099          | 12,8 | 354,938 |
| 2          | 2.777          | 4,197     | 10           | 0.711 |                | 12,647,531                | ) | 4.642          | 46,027     | 10 | 0.082          | 13,0 | 080,690 |
| 3          | 2.098          | 51,689    | 11           | 0.595 |                | 16,133,250                | 3 | 2.601          | 1,853,531  | 11 | 0.072          | 13,2 | 252,088 |
| 4          | 1.668          | 299,677   | 12           | 0.493 |                | 19,598,218                | 4 | 1.190          | 6,191,411  | 12 | 0.061          | 13,4 | 401,719 |
| 5          | 1.430          | 951,177   | 13           | 0.406 |                | 22,900,151                | 5 | 0.566          | 9,396,976  | 13 | 0.056          | 13,5 | 531,668 |
| 6          | 1.264          | 2,133,567 | <b>*</b> 33_ | 0.025 |                | 43,852,665                | 6 | 0.808          | 11,107,361 | 49 | 0.001          | 14,9 | 939,845 |
| 7          | 1.118          | 3,931,575 | 50           | 0.011 |                | 4 <del>6,0</del> 75,896   | 7 | 0.187          | 12,015,748 | 50 | 0.001          | 14,9 | 946,730 |

A low-order modeler is usually unable to capture the correlations between consecutive characters in the text By switching to higherorder models better statistics can be obtained, but the number of different contexts increases, consuming more space. The average length of a word is around 5 bytes in English texts, but the variance is relatively high. In general, a high-order modeler needs to use <u>*k* around 10</u> to capture the correlation between <u>2 consecutive words</u>.



#### Text approx.50 Mbytes

|   | Plain Text     |           |    |       |            |   | Text compressed with ETDC |            |    |                |            |  |
|---|----------------|-----------|----|-------|------------|---|---------------------------|------------|----|----------------|------------|--|
| k | H <sub>k</sub> | contexts  | k  | $H_k$ | contexts   | k | H <sub>k</sub>            | contexts   | к  | H <sub>k</sub> | contexts   |  |
| 0 | 4.888          | 1         | 8  | 0.972 | 6,345,025  | 0 | 7.137                     | 1          | 8  | 0.132          | 12,531,512 |  |
| 1 | 3.591          | 96        | 9  | 0.837 | 9,312,075  | 1 | 6.190                     | 256        | 9  | 0.099          | 12,854,938 |  |
| 2 | 2.777          | 4,197     | 10 | 0.711 | 12,647,531 | 2 | 4.642                     | 46,027     | 10 | 0.082          | 13,080,690 |  |
| 3 | 2.098          | 51,689    | 11 | 0.595 | 16,133,250 | 3 | 2.601                     | 1,853,531  | 11 | 0.072          | 13,252,088 |  |
| 4 | 1.668          | 299,677   | 12 | 0.493 | 19,598,218 | 4 | 1.190                     | 6,191,411  | 12 | 0.061          | 13,401,719 |  |
| 5 | 1.430          | 951,177   | 13 | 0.406 | 22,900,151 | 5 | 0.566                     | 9,396,976  | 13 | 0.056          | 13,531,668 |  |
| 6 | 1.264          | 2,133,567 | 33 | 0.025 | 43,852,665 | 6 | 0.308                     | 11,107,361 | 49 | 0.001          | 14,939,845 |  |
| 7 | 1.118          | 3,931,575 | 50 | 0.011 | 46,075,896 | 7 | 0.187                     | 12,015,748 | 50 | 0.001          | 14,946,730 |  |

The average code length in ETDC is less than **2 bytes**, and the variance is low, as codes rarely contain more than 3 bytes. Hence a *k-modeler* can capture correlations between consecutive words with a **much smaller K**.

- However  $H_k$  values are not directly comparable.
  - ETDC (T) has approx. 1/3 of the symbols of T.
  - Compressors do not use a fixed k, but rather administer in the best way they can a given amount of memory to store contexts.
  - The correct comparison is between the entropy achieved as a function of the number of contexts necessary to achieve it.













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Indexed text



- SSA (Succinct Suffix Array)
  - V. Mäkinen & G. Navarro.
  - Obtains a size related to  $H_0$
- AF-FMindex (Alphabet-Friendly FM-index)
  - P. Ferragina, G. Manzini, V. Mäkinen & G. Navarro
  - Compression approaches  $H_{k.}$
  - We expect <u>AF-FMindex</u> to be successful in detecting high-order correlations in TH(T), where a **smaller** k would be sufficient to succeed compared to that built on T.
  - Important because the AF-FMindex is limited in practice to obtain entropies of relatively low k value.



- Self-indexes' are able to:
  - *Count* the number of ocurrences of a pattern p in O(|p|) steps.
  - *Locate* the position of a suffix in the text.
  - Recover the original text (display / extract).



TH generates suffix-free codes  $\rightarrow$  no false matchings occur



- Self-indexes' are able to:
  - Count the number of ocurrences of a pattern p in O(|p|) steps.
  - Locate the position of a suffix in the text.
  - Recover the original text.
- We chose TH as the base compressor because it generates suffix-free codes.
  - This permit to compress the searched pattern *p* and then search for its compressed form directly.
  - As those self indexes use a terminator (\$) for the indexed text, we modified *TH* to ensure that at least 1 byte value does not appear in the compressed text.





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### **Experimental results**

| STOLLO1 |
|---------|
|         |
|         |
|         |
|         |

| <u>CORPUS</u> | <u>size (bytes)</u> | <u>Num. words</u> | <u>Nº different words</u> |
|---------------|---------------------|-------------------|---------------------------|
| CALGARY       | 2,131,045           | 528,611           | 30,995                    |
| FT91          | 14,749,355          | 3,135,383         | 75,681                    |
| CR            | 51,085,545          | 10,230,907        | 117,713                   |
| FT92          | 175,449,235         | 36,803,204        | 284,892                   |
| ZIFF          | 185,220,215         | 40,866,492        | 237,622                   |
| FT93          | 197,586,294         | 42,063,804        | 291,427                   |
| FT94          | 203,783,923         | 43,335,126        | 295,018                   |
| AP            | 250,714,271         | 53,349,620        | 269,141                   |
| ALL FT        | 591,568,807         | 124,971,944       | 577,352                   |
| ALL           | 1,080,719,883       | 229,596,845       | 886,190                   |

- Intel Pentium-IV 3 Ghz 4Gb RAM.
  - Debian GNU/<u>Linux</u> (kernel 2.4.27)
  - gcc 3.3.5 and optimization <u>–O9</u>
  - Time measures <u>CPU user-time</u>

# Experimental results Compression ratio









# Experimental results Compression time







In seconds

## Experimental results Decompression time





In seconds

# Experimental results Compression ratio





Comparison against other ppm-based algorithms that use a high value of k.

Monstruous ppmETDC +(ppmd var J)ppm-monst(k=128)15.76%

All the co-occurrences of the symbols in the text have been detected.

### Experimental results Indexing



•We used the corpus CR (aprox 50 Mbytes)

Size of index: Compression ratio (%)

| Rank Factor 16 | Sample Rate |        |        |        |  |  |  |
|----------------|-------------|--------|--------|--------|--|--|--|
|                | 16          | 32     | 64     | 1024   |  |  |  |
| TH + affm      | 49,95%      | 41,84% | 37,78% | 34,73% |  |  |  |
| Plain + affm   | 104,83%     | 79,83% | 67,33% | 57,96% |  |  |  |
| TH + ssa       | 47,94%      | 43,88% | 41,86% | 40,33% |  |  |  |
| Plain + ssa    | 111,69%     | 99,19% | 92,94% | 88,25% |  |  |  |
| тн             | 34,31%      |        |        |        |  |  |  |

### Experimental results Indexing





### By setting SR=1024 and RF=64...

- —Less space (but slower indexes at searches)
  - the AF-FMindex occupies less than the text compressed with TH

| Compression ratio |        |  |  |  |
|-------------------|--------|--|--|--|
| TH+affm           | 32.71% |  |  |  |
| Plain+affm        | 53.59% |  |  |  |
| TH+ssa            | 38.37% |  |  |  |
| Plain+ssa         | 83.62% |  |  |  |
| ТН                | 34.31% |  |  |  |

### Experimental results Indexing



- For each SR value there is a line depending on the **RF** values
- We measured time in ms (for locate).



### Outline





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# Conclusions





- By preprocessing a text T with either ETDC or TH:
  - We obtain a compressed text of around 30% of size(T).
  - Still compressible and indexable.
- By compressing in a <u>second step</u> with PPM, gzip or bzip2, we improve: compression ratio, compression speed, and decompression speed.
  - ETDC+gzip: very fast and good compression ratio (<bzip2)</p>
  - ETDC+bzip2 compresses a little bit more, at the expense of a lower speed.
  - ETDC+PPM: the best compression (but still slow).

# Conclusions





If we apply the AF-FMindex and SSA over text compressed with TH and:

- we set the index parameters in order to obtain a structure of the same size as if we indexed the plain text, we obtain two indexes that are much faster than the traditional ones.
- If we instead set the parameters to obtain two indexes with the same search speed, the index over the compressed text will occupy around 30% less than in the case of the plain text.